FUTURES IMPLEMENTATION IN XINU

Introduction:

The design document provides an overview of the futures concept part of the new C++ standard. The futures variable is a placeholder for an eventual output generated by a function called asynchronously.

The asynchronous function can then use a variety of methods to query, wait for, or extract a value from the future variable. These methods may block if the asynchronous operation has not yet provided a value.

Design Considerations

- The program has producers and consumers of data. The producer outputs a value and sets the future variable using future_set() and the consumer consumes data and it queries the future variable using the future_get() as per our implementation.
- If a thread calls future_get() on an empty future, then the calling thread should block and subsequent future_get() calls should fail. If a thread calls future_set() on an empty future, then it becomes full and a subsequent future_set() should fail. Calling future_get() on a full future yields the value and resets the future's state to empty.

Implementation Details

- Header file <future.h> contains function prototype and the structure corresponding to the future variable.

```
struct futent          /*future table entry*/
{
        int data;          /* Data that future should hold */
        char state;   /*state of the future which can be FUT_FREE or FUT_USED*/
        int flag;          /*flag to check if a process is waiting for the future*/
        int tid;           /*process id waiting for the future*/
};
```

- future futalloc() in futcreate.c

  - The function futalloc () returns the ID of the allocated future to f1.
    future f1;
    f1=futalloc();

- syscall future_set(future fut, int value)
    - The function is implemented as a system call and it is used to set the data of the future.
- syscall future_get(future fut, int *value)
    - The function is implemented as a system call and it is used to get the data stored in future.
- syscall future_free(future fut)
    - The function is implemented as a system call to free up the future variable.

Input snapshot

```
int a,b; a=10; b=20;
future f1,f2;
f1=futalloc();
f2=futalloc();
resume(create(consumer,1024,20,"cons1",1,f1));
resume(create(consumer,1024,20,"cons3",1,f1));
resume(create(producer,1024,20,"prod1",2,f1,a));
resume(create(producer,1024,20,"prod2",2,f2,b));
resume(create(producer,1024,20,"prod3",2,f2,b));
resume(create(consumer,1024,20,"cons2",1,f2));
```

Output snapshot

```
Output: Trying to access blocked future: Process cons3 blocked
Calling used future, Process prod3 blocked
The value is 10
The value is 20
```

• cons1 process executes and waits for the future f1 value and is invoked once future f1 value is set.
• cons3 process gets blocked since cons1 process is already waiting for future f1.
• cons2 process executes and waits for future f2 since the value is not set initially and resumes execution once future f2 value is set.
•prod3 process gets blocked since future f2 is already used by process prod2 and is full.

References

Wikipedia, CPP reference, Xinu Approach- Douglas Comer